Reading: Patt & Patel

Chp 6.1.2-6.1.3 (control constructs and mechanisms)

Chp 6.2.1 up through Example 1. (Debugging, using the simulator/debugger.

Chp. 7.1-7.4 (LC-3 Assembly language: instruction syntax, labels, comments, assembler directives, 2-pass assembly, object files and linking, executable images).

Also, see in docs/:

LC3-assemblySyntax.html LC3-IntructionsSummary.html README-LC3tools.html LC3-assemblyCheatSheet.asm LC3-AssemblyManualAndExamples.pdf

Problems:

P&P, Chp. 6 6.10 (conditional, odd?) 6.12a (char echo) 6.15 (figure out instruction)

P&P, Chp. 7

7.1 (hand assemble) NB--Assume we mean "the location in the object file that will get loaded to memory location x3025" for the text's "location x3025 of the object file."

7.2 (hand assembly, symbol table entry)

7.4 (hand assembly, symbol table)

7.8 (hand assembly, program trace)

7.14 (debugging asm)

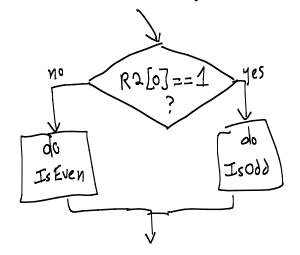
7.24 (looping control)

7.25 (too large constant in .FILL)

NB--Using a simulator/debugger in conjunction with an assembler can be helpful for these problems. Assembler is lc3as; simulator/debugger is PennSim.jar.

Chp 6/7 Solhs

6.10) Content of R2 even or odd? Give machine language. Use conditional construct.



; R2 even or odd ; If odd, bit_0 is 1	; else even.
AND R0, R2, BRz IsEven	1 ; result depends only on bit_0
IsOdd: jmp Done	; do odd stuff
IsEven: Done:	; do even stuff

(6.12) Echo one keyboard (kb) char to display.

,_____

;---- KBecho

;---- Get char from kb, send to display. We take the lazy route and use pre-coded OS functions.

;---- That's what they are there for, after all. "GETC" and "OUT" are assembler aliases

;---- for TRAP x20 and TRAP x21, respectively.

TRAP x20 ;---- "GETC": Read char from kb into R0[7:0]

TRAP x21 ;---- "OUT": Send char in R0[7:0] to display.

GETC does polling until the kb controller signals it has new data from the keyboard. GETC then reads the content of kb's data register into R0. OUT polls until the display signals it can accept new data, then writes the content of R0 to the display's data register. Note that kb zeroes the high eight bits of its data register and puts an ASCII code in the low eight bits; the display reads only the low eight bits of its data register.

To poll a status register's ready bit (bit_15), read the device's status register into a register, then check whether the register's high bit is 1. We access kb's status register via memory address xFE00; its data register via memory address xFE02. For code clarity, we keep those addresses in memory words in our program:

KBSR:	.FILL xFE00
KBDR:	.FILL xFE02

and reference them as pointer variables via their labels, eg., LDI R0, KBSR

For code details of OUT and GETC, see Ic3os code:

(1) using an LC3 debugger/simulator, find the addresses of the two trap routines in the trap vector table, then look in those memory locations to find the code. The trap vector for GETC is in memory location x0020; so, look at memory location x0020, write down the address found there, and then look at memory content at that address. Vector for OUT is at x0021. Realize that the code is in memory because the debugger/simulator loaded it to memory. That is something that would happen as part of system start-up, called "boot strapping" or booting for short;

(2) or see OS source code in src/lc3os.asm or versions thereof.

(6.15) Given reg and mem content, what instruction @ x3010 was executed?

There were no changes in reg content, only mem x3406 changed from x31BA to xE373. R2 contained xE373. Therefore, the instruction must have been a memory write of R2 via ST or STI or STR.

We try ST: ST R2, <offset> We know the address must come out to be x3406 and the PC contained x3010+1; so x3011 + <offset> = x3406

We get $\langle offset \rangle = x3F5$, positive 2's complement, or 011 1111 0101 in as few bits as possible. That's 11 bits, but the offset field for ST is only 9 bits. ST won't work. We could try STI, but there isn't any memory location containing the address x3406; so, there's no point.

It must be STR: STR R2, R?, <offset>

STR's address arithmetic says:

R? + <offset> = x3406

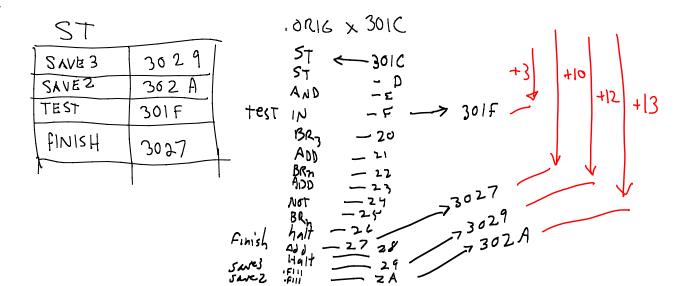
where the offset is 6-bit 2's complement. The register value closest to x3406 is in R4, x33FF. The difference is +7. So, the offset is x07, or 00 0111 in 6 bits.

STR R2, R4, #7

(7.1) What's in x3025 after assembly?	; in effect, code is, .orig x3024 place: .FILL x45A7 LDI R3, place	assembly makes memory as: x3024: x45A7 x3025: (bits for "LDI R3, -2")	result is: x3024: 0100010110100111 x3025: 1010 011 111111110 or x3024: x45A7 x3025: A7FE
(7.2) ASCII LD R1, ASCII What's in R1 after execution?	×4F08 : LD RI,	<pre><offset: pc+offset="x4F08"></offset:></pre>	Pc= ×4F09 → offset = -1 ×23ff

(7.4) ST for

code?



(7.8) Show REG_FILE content at breakpoint.
(xA400) this1: lea r0 this1
(xA401) this2: ld r1, this2
(xA402) this3: ldi r2, this5
(xA403) this4: ldr r3, r0, #2
(xA404) this5: .fill xA400
bp = xA404

$$P_{2} = x EIFF$$

$$C_{10} = x A400$$

$$R_{2} = x EIFF$$

$$C_{10} = x A400$$

$$R_{3} = xA401$$

$$R_{4} = 0$$

$$R_{4} = 0$$

$$R_{4} = 0$$

(7.14) Assemble code

(a) , ORIG X3000
 STE RO, LABEL → IOII 000 0 0000 0010 × B002 ①
 OUT → IIII 0000 0010 000 × F021 ②
 HALT → IIII 0000 0010 0101 × F025 ③
 LABEL , STRINGZ "7" → 0000 0000 0000 0000 0000 × 0000 Ø
 , END (c) STORES × 3000 into TVT for thep × 25, choses
 ③ k jmp to ①; ③ outputs arcii NUL to display