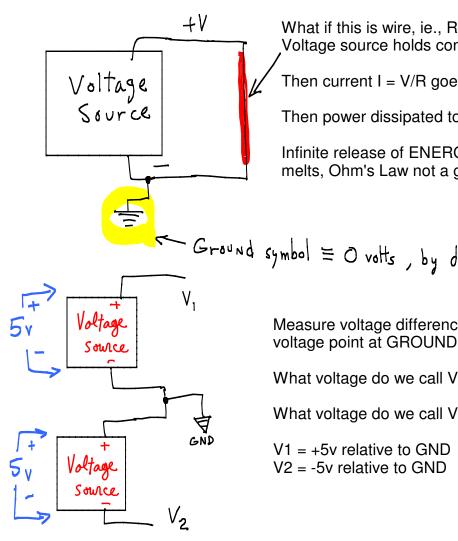


Ohnn's Law, conductors, ...

$$R = \frac{1}{2} V_{I}$$
 (i) relationship is Linear, then ohn's Lea Resister)
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 $V = \frac{1}{2} V_{I}$ (i) (i) relationship is Linear, then ohn's Lea Resister)
 $V = \frac{1}{2} V_{I}$ (i) (i)



What if this is wire, ie., R is nearly 0? Voltage source holds constant V voltage difference.

Then current I = V/R goes to infinity.

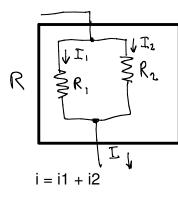
Then power dissipated to heat = I^*V goes to infinity.

Infinite release of ENERGY! BOOM! (Actually, wire and/or supply melts, Ohm's Law not a good approximation at that point.)

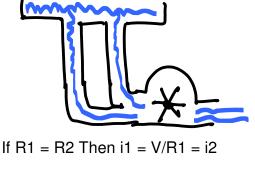
Measure voltage difference from voltage point at GROUND.

What voltage do we call V1?

What voltage do we call V2?



What is total R?

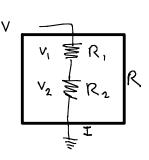


R = V/i = V/(2V/R1) = 1/2 R1

Resistors connected in PARALLEL.

Same as two water pipes in parallel: less resistance to flow, total flow is sum of flow in both.

Water pressure, voltage, is same for both paths.



Resistors connected in SERIES.

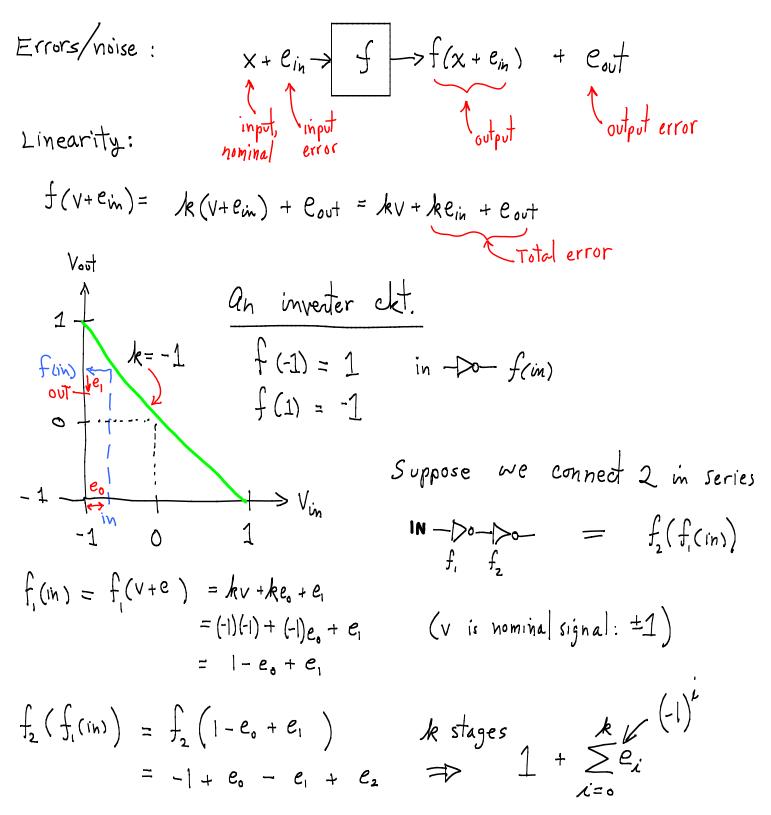
Same as two pipes end-to-end: more resistance, less flow, same flow in both.

$$\begin{array}{rcl} R &=& V/i &=& (V1+V2)/i \\ V1 &=& i^*R1 & V2 &=& i^*R2 \\ R &=& (i^*R1+i^*R2)/i &=& (R1+R2) \end{array}$$



We need Signal-Restoring, Non-Linear Logic. Ohm's Law devices are LINEAR.

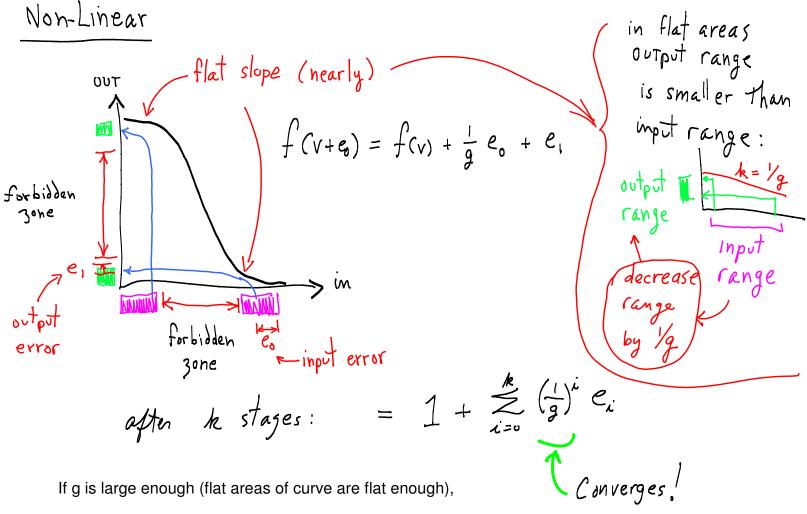
Suppose we had only linear devices (or something very nearly linear), then signal output has errors proportional to input errors.



The errors include signs = random walk with random size steps.

Errors independently random w/ average 0 ==> variance increases w/ k. Total error grows w/o bound! Take random step (either in the -1 or +1 direction). How far from 0 can you expect to be after k steps? About $k^{1/2}$ away. With probability 0 you will be at 0, and error gets unboundedly large.

We must Reduce error at each stage ==> exponentially decreasing effect in later stages.



and

if output error size is not too big,

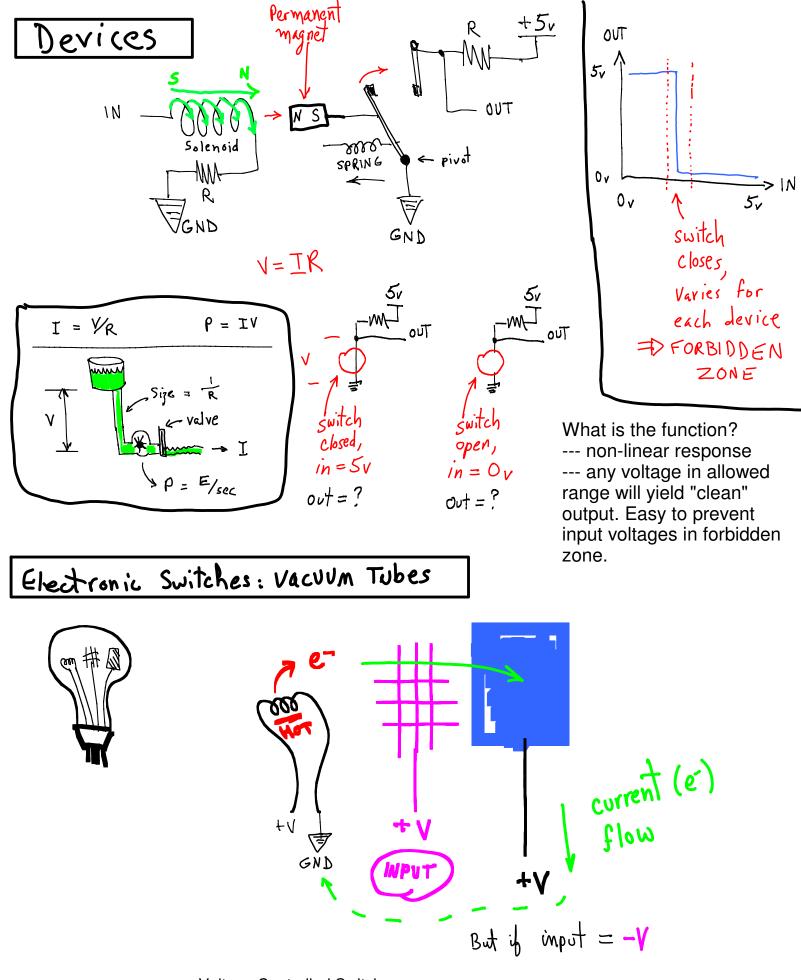
Then

output after k stages never hits FORBIDDEN ZONE.

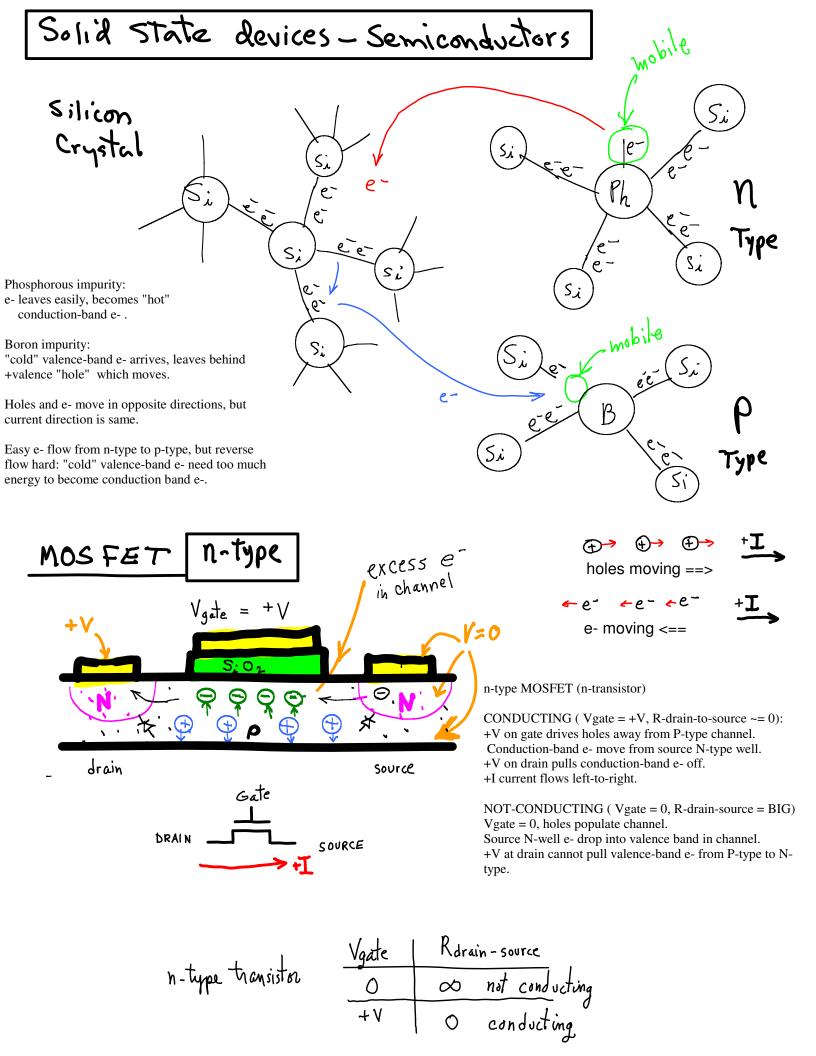
So, if we plan to have a circuit with long device chains, we must have non-linear devices w/ suitable response curves.

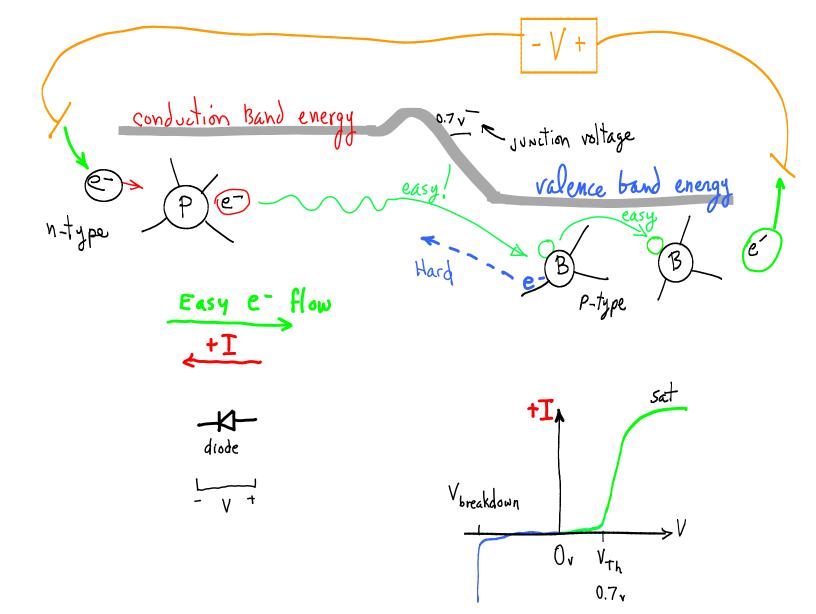
Do we plan to have long chains? YES:

- (1) feedback in system,
- (2) chained data operations: D1 => D2 => D3 => D4 ...
- (3) 1 Billion devices per cpu

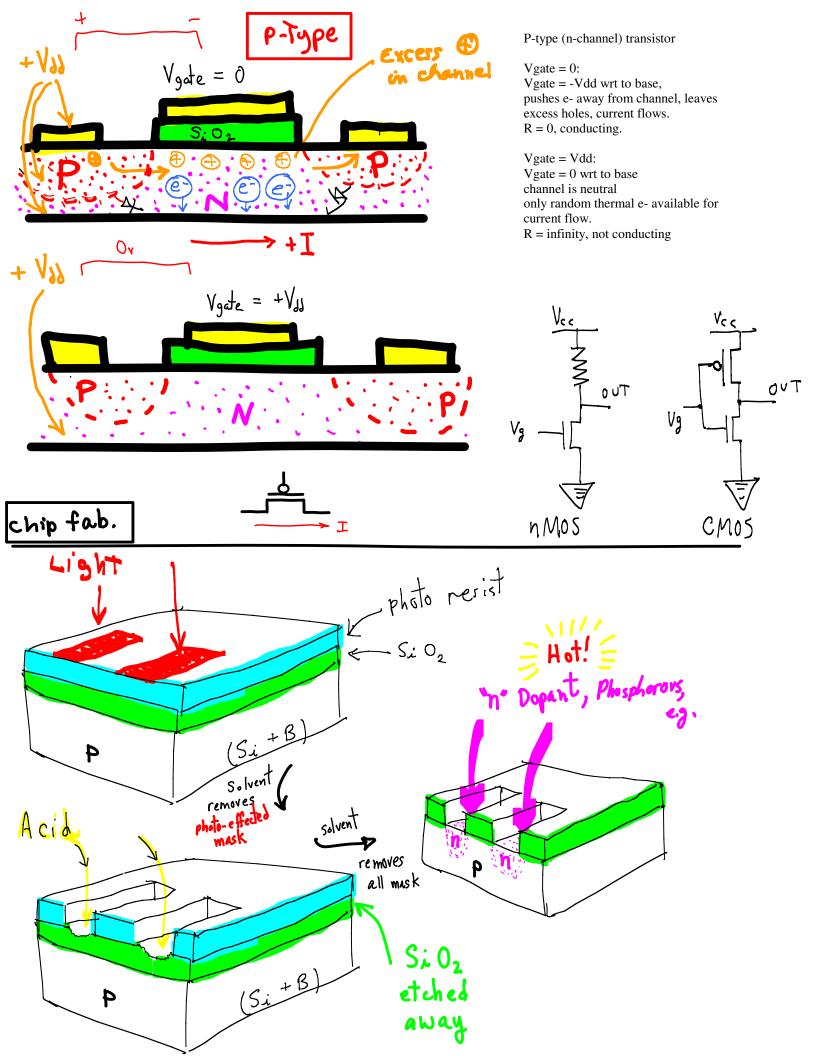


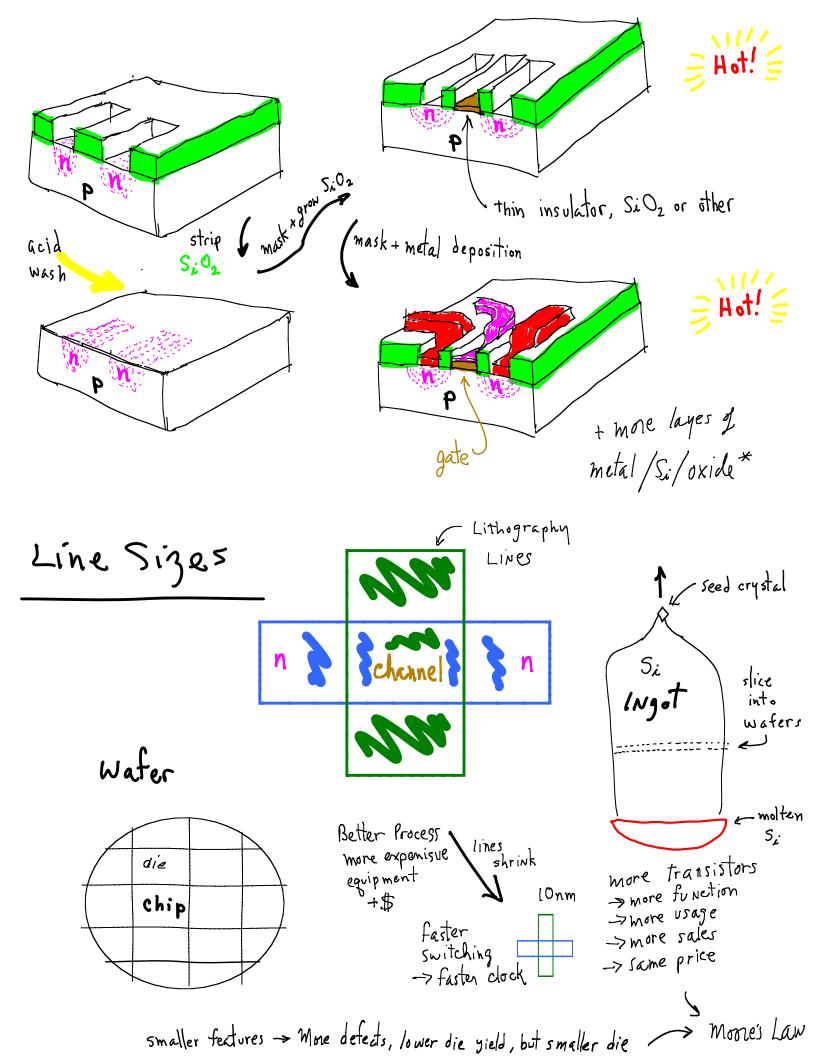
Voltage Controlled Switch faster, more reliable, less power



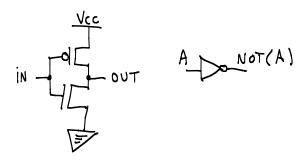


Just what we want: nice non-linear switch.





Basic Logic Gates



Pz

P₁

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Nz

*

- not conducting

 \ast

OUT

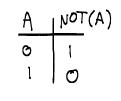
Ö v

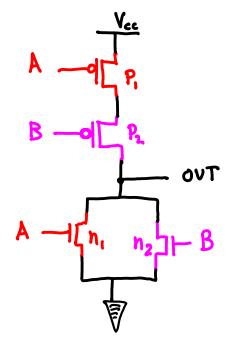
() v

٧دد

Ö v

N,





A B out ------0 0 || 1 0 1 || 0 1 0 || 0 1 1 || 0

NOT (OR)

NOR

Two kinds of logic circuits:

CIRCUITS

Logic

(1) w/ feedback, SEQUENTIAL: can hold STATE

(2) w/o feedback, COMBINATORIAL: realize FUNCTIONS

Basic logic gates: NOT, 2-input NOR, 2-input NAND. That's all we need for both sequential and combinatorial circuits.

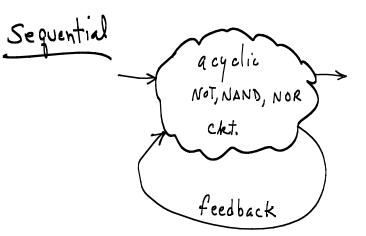
A

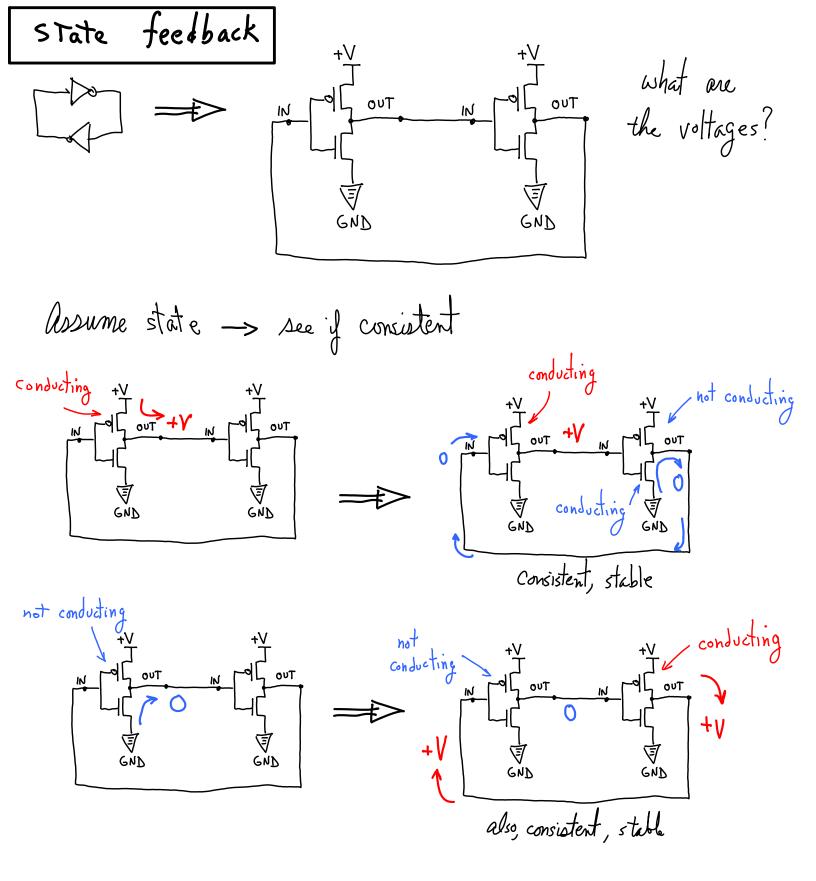
0

B

(NAND alone is sufficient, also NOR alone is sufficient.)

Combinatoria acyclic NOT, NOR, NAND ckt

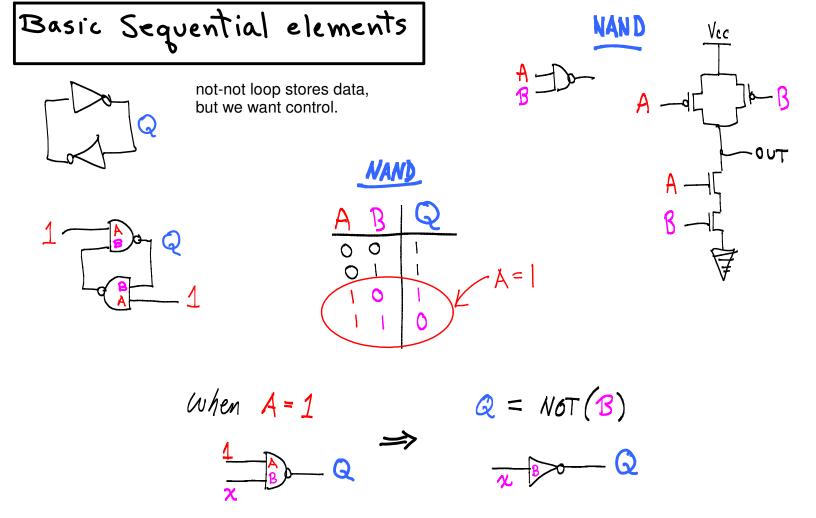




Circuit has two stable states: aka, meta-stable, bi-stable.

What state when power is first turned on? Unknown, random.

Can we set the state, using voltage inputs? No, useless!?!



CIRCUITS W/ STATE

NOT-NOT circuit is stable in either of two states: BISTABLE element.

NAND-NAND circuit with both A inputs = 1, same as a NOT-NOT circuit.

