Lisp!

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Why Lisp?

- Great for symbolic computation
- Great for functional programming
- Learning Lisp will make you a better programmer in other languages

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- Phillip Greenspun's Tenth Rule of Programming: "Any sufficiently complicated C or Fortran program contains an ad hoc, informally-specified bug-ridden slow implementation of half of Common Lisp."

Symbolic Computation

- Differentiation:
 - ▶ General form: $a^d \Rightarrow da^{d-1}$
 - ▶ Application: $n^2 \Rightarrow 2n$
- DeMorgan's Law:
 - ▶ General form: $\neg(\phi \land \psi) \Rightarrow \neg \phi \lor \neg \psi$
 - ▶ Application: $\neg(\neg p \land q \rightarrow r) \Rightarrow \neg \neg p \lor \neg(q \rightarrow r)$

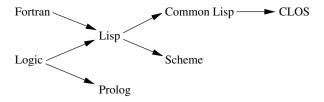
Functional Programming Languages

- Everything is a function that returns a value, even for loops
- ► First-class and higher-order functions
- Purely functional languages have immutable data structures, which eliminates side effects
- ▶ No side effects means computation can be distributed
- Pure functional: Haskell
- "Impure" functional: Common Lisp, Clojure
- Languages have different degrees of purity.

What is Lisp?

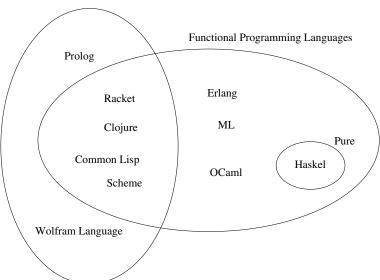
- Specified by John McCarthy in 1958
- ► LISP ≡ LISt Processing
- Functional programming language
- Great for symbolic computing; important for AI tasks of the day
 - Predominant languages of the day were Fortran and COBOL
- Interpreted, but it can be compiled
- Great for rapid prototyping

A Lineage of Lisp



Languages for Symbolic Computing

Languages for Symbolic Computing



Why Aren't We Learning Clojure?

- Clojure is a dialict of Lisp that runs on the Java Virtual Machine
- Immutable data structures
- Optimized for concurrency via map-reduce
- There are oddities (e.g., doseq returns nil)
- Clojure is still developing
- My advice: Go learn it after this class

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- Functions, function calls, and control structures are also linked lists

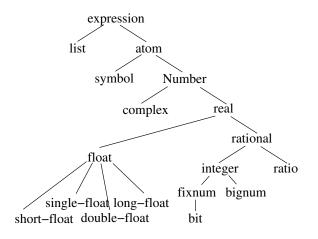
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- And then there are the parentheses...

Lisp Data Types

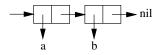


Lists and Function Calls

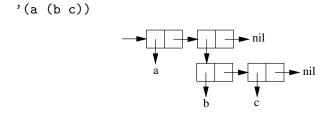
- ▶ We form lists and function calls using parentheses
- Prefix notation:
 - ▶ The first element of a list is the operator or function name
- Use a single quote to prevent evaluation
 - ► Empty list: ()
 - ▶ List: '(sqrt 4)
 - ► Function Call: (sqrt 4)
 - ▶ If-statement: (if pig 'cow 'dog)
 - ▶ List: '(if pig 'cow 'dog)

Linked Lists in Lisp

'(a b)



Linked Lists in Lisp



Numeric Operators and Functions

- _ _
- . .
- **>** *
- ▶ sqrt
- ▶ 1+
- ▶ 1-
- ► min
- ▶ max
- ► expt
- ▶ log
- ▶ abs
- ▶ sin, cos, tan
- asin, acos, atan
- ▶ floor, ceiling, round, truncate

Output Functions

- ▶ (format ⟨file-stream⟩ ⟨control-string⟩ ⟨variable-list⟩)
- ► Example: (format t "This is a test, ~a~%", 'bob)
- Output: This is a test, BOB
- Basic control codes:
 - ▶ ~a any lisp object
 - ▶ ~s s-expression
 - [™]d decimal
 - ▶ ~f float
 - ► ~% newline character
- Note that in sbcl ouptut is buffered by default.
- ▶ (flush-output).

prog1 Blocks

```
Syntax:

\begin{array}{c} (\operatorname{prog1} \\ \langle \mathit{form}_1 \rangle \\ \langle \mathit{form}_2 \rangle \\ \dots \\ \langle \mathit{form}_n \rangle) \end{array}
```

- Semantics:
 - ▶ Evaluates forms 1–n
 - ▶ prog1 evaluates to the valuation of ⟨form₁⟩

progn Blocks

```
Syntax:

\begin{array}{c} (\text{progn} \\ \langle \textit{form}_1 \rangle \\ \langle \textit{form}_2 \rangle \\ \dots \\ \langle \textit{form}_n \rangle) \end{array}
```

Semantics:

- ▶ Evaluates forms 1–n
- progn evaluates to the valuation of $\langle form_n \rangle$

Defining Functions: defun

```
(defun \langle function-name \rangle (\langle parameter-list \rangle) \langle form \rangle)
```

- Semantics:
 - ▶ Defines a function with the name ⟨function-name⟩ with the parameters ⟨parameter-list⟩ and the body ⟨form⟩.
 - Evaluates to the function \(\langle function-name \rangle \)

Binding Local Variables: let and let*

```
 \begin{array}{c} (\mathsf{let}[*] \, ((\langle \mathit{variable}_1 \rangle \, \, \langle \mathit{value}_1 \rangle) \\ \hspace{0.5cm} (\langle \mathit{variable}_2 \rangle \, \, \langle \mathit{value}_2 \rangle) \\ \hspace{0.5cm} \dots \\ \hspace{0.5cm} (\langle \mathit{variable}_n \rangle \, \, \langle \mathit{value}_n \rangle)) \\ \hspace{0.5cm} \langle \mathit{form} \rangle) \end{array}
```

- Semantics:
 - ▶ Defines locally scoped variables 1−n
 - ▶ Binds or assigns values 1-n to variables 1-n, respectively
 - ► Evaluates ⟨*form*⟩
 - ▶ let evaluates to the valuation of ⟨form⟩
 - ▶ let* forces the sequential assignment of variables 1-n

Branching: if-then-else, when, unless

Syntax:

```
(if \langle test \rangle \ \langle then	ext{-}form \rangle \ [\langle else	ext{-}form \rangle])
```

Semantics:

- ► Evaluates ⟨*test*⟩
- ▶ If ⟨test⟩ is not nil, evaluates ⟨then-form⟩
- ▶ Otherwise, if ⟨*else-form*⟩ is not nil, evaluates ⟨*else-form*⟩
- if evaluates to the form evaluated

Also:

- (when $\langle test \rangle \langle form \rangle$) \equiv (if $\langle test \rangle \langle then-form \rangle$)
- (unless $\langle test \rangle \langle form \rangle$) \equiv (if $\langle test \rangle$ nil $\langle else\text{-}form \rangle$)
- (unless (not $\langle test \rangle$) $\langle form \rangle$) \equiv (if (not $\langle test \rangle$) $\langle then\text{-}form \rangle$)

Relational and Equality Operators and Functions

- equalp same expression?
- equal same expression?
- eq1 same symbol or number?
- ▶ eq same symbol?
- ► = same number?
- ▶ atom is it an atom?
- numberp is it a number?
- consp is it a cons?
- listp is it a list?
- symbolp is it a symbol?
- boundp is it a bound symbol?
- Also member, null, zerop, plusp, minusp, evenp, oddp, <, <=, >=, and >.

Nested if-then-else Statements: cond

```
(\operatorname{cond}(\langle \operatorname{test}_1 \rangle \langle \operatorname{consequent}_1 \rangle) 
(\langle \operatorname{test}_2 \rangle \langle \operatorname{consequent}_2 \rangle) 
\dots
(\langle \operatorname{test}_n \rangle \langle \operatorname{consequent}_n \rangle))
```

- Semantics:
 - ▶ Evaluates tests 1–n
 - Evaluates the consequent of the first test that is not nil
 - cond evaluates to the valuation of the consequent or to nil if no test is not nil
 - ▶ Note that the literal t always evaluates to true

Looping through Lists: dolist

```
 \begin{array}{c} (\texttt{dolist} \ (\langle \textit{variable} \rangle \ \langle \textit{list-form} \rangle \ [\langle \textit{result-form} \rangle]) \\ \langle \textit{form} \rangle) \end{array}
```

- Semantics:
 - ▶ Iterates once for each element in ⟨*list-form*⟩
 - ► Assigns the element to ⟨variable⟩
 - ► Evaluates ⟨*form*⟩
 - ▶ Upon loop termination, evaluates ⟨result-form⟩
 - ▶ dolist evaluates to the valuation of ⟨result-form⟩

Count-controlled Loops: dotimes

```
 \begin{array}{c} (\texttt{dotimes} \ (\langle \textit{variable} \rangle \ \langle \textit{upper-bound-form} \rangle \ [\langle \textit{result-form} \rangle]) \\ \langle \textit{form} \rangle) \end{array}
```

- Semantics:
 - ► Iterates over [0, ⟨upper-bound-form⟩)
 - ► Assigns the counter value to ⟨variable⟩
 - ► Evaluates ⟨*form*⟩
 - ▶ Upon loop termination, evaluates ⟨result-form⟩
 - ▶ dotimes evaluates to the valuation of ⟨result-form⟩

General Looping Construct: do and do*

Syntax:

```
 \begin{split} (\text{do}[*]((\langle \textit{var}_1 \rangle \ \langle \textit{init}_1 \rangle \ \langle \textit{update}_1 \rangle) \\ & (\langle \textit{var}_2 \rangle \ \langle \textit{init}_2 \rangle \ \langle \textit{update}_2 \rangle) \\ & \cdots \\ & (\langle \textit{var}_n \rangle \ \langle \textit{init}_n \rangle \ \langle \textit{update}_n \rangle)) \\ & (\langle \textit{termination-test} \rangle \ \langle \textit{result-form} \rangle) \\ & \langle \textit{form} \rangle) \end{split}
```

Semantics:

- Defines locally scoped variables (var_i)
- ► Initializes them to ⟨init_i⟩
- ▶ Loops until ⟨termination-test⟩ is not nil by
 - updating $\langle var_i \rangle$ by evaluating $\langle update_i \rangle$
 - evaluating \(\form \rangle \)
- ▶ Upon loop termination, evaluates ⟨result-form⟩
- ▶ do evaluates to the valuation of ⟨result-form⟩
- do* forces the sequential assignment of \(\lambda var_i \rangle \)

Arrays

```
Creation
```

```
(setf single (make-array '(4)))
```

```
▶ (setf double (make-array '(2 4)))
```

•

```
(setf double (make-array '(2 4) :initial-element nil))
```

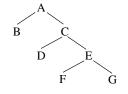
Access

- ▶ (format t "~a~%" (aref single 1))
- ▶ (format t "~a~%" (aref double 0 0))

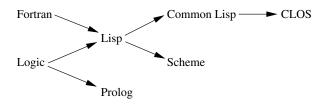
Assignment

- ▶ (setf (aref single 1) 'hoya)
- ▶ (setf (aref double 0 0) 'saxa)

Binary Trees



Graphs



```
(setf (get 'fortran 'influenced) '(lisp))
(setf (get 'logic 'influenced) '(lisp prolog))
(setf (get 'lisp 'influenced) '(scheme common-lisp))
(setf (get 'common-lisp 'influenced) '(clos))
```

Things We Won't Cover

- Dotted Lists
- Types
- Strings
- Serious File I/O
- Packages
- Hash Tables
- Structures (aka records)
- Exceptions: throw and catch
- Programmer-defined macros
- ► CLOS: classes, properties, methods
- Many, many functions, forms, and macros