# Algorithms for NLP

#### **Finite State Morphology**

Amir Zeldes <u>amir.zeldes@georgetown.edu</u>

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#### Regular Expressions – a language?

- By now we've all seen regular expressions
   Very useful for finding phone numbers, URLs
   Or potentially
  - catching killers:





#### Regular Expressions – a language?

Can regex capture the grammar of a language?
What is the grammatical structure of e-mail?





## From regex to natural language

- Regular expressions describe a simple grammar
  - For example, you could think of a regex: /DA\*N/
  - As describing a Noun Phrase:

Determiner	(Adj)*	Noun
The	quick brown	fox
a		house
ту	lovely	cat

- Just replace each noun with N, each Adj with A...
- We can now recognize noun phrases! (why should we?)

## From regex to natural language

 In fact, syntax is more complex than what we can express with regex:

pick the kids up: /VDNP/

- But only certain verbs take certain particles, objects...
- Can't prevent matching:
   sleep the kids up
   pick the kids over



## From regex to natural language

- But for morphology, word formation is often describable using something like regex:
  - Super anti adverbs: /(super)?(anti)?ADJ(ly)?/ super-anti-ingenious-ly
  - Noun compounds: /N+N/ nightgown
- But what is ADJ? or N?
- Can we do regex with a different 'alphabet'?
- A grammar of expressions using any 'alphabet' is called a regular language



# Regular languages

- In fact we can create a regular language grammar using:
  - Some alphabet Σ with symbols a, b, c...
  - Any single symbol is a possible regular grammar (just a)
  - Any union (a OR b), concatenation (a THEN b) or Kleene star (a\*) of a symbol or language
- Using these constraints, we can build any regular grammar using any set of symbols



#### Finite State Methods

 Another way to look at regular grammars is thinking of a reader head

0 0 0 0

0

q

1 1

В

0

- Moving from character to character on a ribbon
- /ba+\$/ matches like this:
  - Initial state: read till you see a b
  - The letter **b** is reached -> change to state 2
  - Move right on the ribbon look for a
    - If **a** is seen -> stay in the same state, keep going
    - Else if non-a is seen -> match fails
    - Else if input runs out -> done (successful match)



#### Finite state automaton - FSA

- This type of computing characterizes an FSA:
  - Finite number of states, including start and end
  - Transitions depend on input
- More formally:
  - FSA = {Q, q0, F,  $\Sigma$ ,  $\delta(q,i)$ }
- Where:
  - Q is a set of possible states qi... qn
  - q0 is the starting state within Q
  - F is a subset of end states within Q
  - Σ is the alphabet
  - δ(q,i) is a set of allowable transitions from state q given input i



## Finite State Morphology

- Among most successful applications of FSA
- Popular for agglutinative languages (Turkish, Japanese), and highly inflected concatenative ones (e.g. Slavic)
- Some approaches to non-concatenative morphologies (Arabic, Hebrew)
- Basic tasks:
  - Morphological parsing
  - Generation



#### From NL input to states

- Famous Turkish example (Jurafsky & Martin 2008, after Kemal Oflazer):
  - UygarlaştIramadIklarImIzdanmIşsInizcasIna civil-bec-caus-npot-part-pl-p1pl-abl-past-2pl-adv "such that you can't be made civilized by us" (civil-ize-ate-unable-ing-s-our-from-did-you-ly)
- Morphemes follow a particular order
- Many are optional
- Possible word formations can be described via states...

# Morphological parsing

- The task:
  - Given some word in language X as input:
  - Output lexicon forms of constituent parts ("morphemes")
  - Give morphological analysis to the units
- Ambiguity is possible:
  - friendly (ADJ) = friend:N + ly:ADJ
  - friendly (ADV) = friend:N + ly:ADJ + 0:ADV (for ?friendlyly, Bauer 1992)
- In ambiguous cases: give all possible analyses



## **English adjectives**

- What would we need to model forms like these?
  - happy, happier, unhappy, happily, unhappily
  - lucky, luckiest, unlucky, luckily, unluckily
  - big, bigger, biggest
- What is the alphabet like?
- What transitions are possible?



#### First approximation

- A first approach would be to model states for each morpheme
- Allow transitions based on order (Antworth 1990)



#### Problems

- Some ungrammatical forms will be possible:
  - bigly
  - Unbiggest
  - ...
- Orthography would need to be handled:
  - happyer
  - happyly



#### Solutions

 Automata must become more complex to model the phenomenon
 Just the beginning:



#### Writing automata

- Many frameworks exist for FSM
- Influential early framework: Xerox FSM (XFSM)
  - Beesley & Karttunen (2003)
- Many (re)implementations:
  - HFSM, Foma, OpenFST/PyFST
  - Compiled in C++ for performance
  - Bindings for Python available (though may be tricky to compile, OS dependent)
  - We will work with Foma today (platform independent)



# **Running Foma**

You should have Foma for your OS from here:

- <u>https://code.google.com/p/foma/</u>
- To run it, open a terminal window
  - Windows: Window key -> cmd
  - Mac: run Terminal
  - Navigate to directory: cd YOUR\_PATH
  - Run foma: foma (or foma.exe)



# **Running Foma**

• Interactive mode is like an interpreter (e.g. in Python): foma[0]: define Consonant [p|t|k|b|d|g|||m|n|r|s|z]; defined Consonant: 543 bytes. 2 states, 12 arcs, 12 paths. foma[0]: define Vowel [a|e|i|o|u]; defined Vowel: 333 bytes. 2 states, 5 arcs, 5 paths. foma[0]: regex Consonant Vowel+; 757 bytes. 3 states, 22 arcs, Cyclic. foma[1]: words # Outputs possible inputs which terminate pa paa pae pe

pea



#### Some regex differences to POSIX

The regex syntax used in Python is called POSIX
XFST / Foma syntax has some differences:

- Use spaces to delimit symbols:
  - cat = a single symbol, called cat
  - c a t = three symbols, called c, a, t
  - {cat} = alternative spelling of c, a, t
- The dot is replace by ?:
  - {c ? t} = cat, cot, cut ...
- Use brackets instead of ? for optional:
  - (UN) ADJ
- % is the escape symbol (like \): %? = a real question mark
- \ is negation: \a = not an a
- The only disjunction (OR) is: [x|y] no ranges like [abc]

## Strings and symbol names

- Definitions are most of the work, look like this:
  - **define NAME** DEFINITION;
- Definitions can contain string literals in {...} or other names:
  - define first\_name {Amir};
  - **define** *last\_name* {Zeldes};
  - define full\_name first\_name last\_name;



## Strings and symbol names

 We can also allow alternative values: define *first* [{Bobby} | {Amir}]; define *last* {Zeldes}; define *full* first last; regex *full*; words;

BobbyZeldes AmirZeldes



# Let's try the English adjectives

- Suppose adjectives look like this:
  - Can start with un-
  - Have a stem like *big* or *clear*
  - Can end in -er, -est
- Use:
  - **define** SYMBOL1 [{string1}|{string2}|...];
  - regex (SYMBOL1) SYMBOL2 ...;
  - words;



## Solution

foma[0]: **define** UN {un}; defined UN: 212 bytes. 3 states, 2 arcs, 1 path. foma[0]: define STEM [{big}|{clear}]; defined STEM: 423 bytes. 8 states, 8 arcs, 2 paths. foma[0]: define SUFFIX [{er}] {est}]; defined SUFFIX: 303 bytes. 4 states, 4 arcs, 2 paths. foma[0]: regex (UN) STEM (SUFFIX); 607 bytes. 13 states, 16 arcs, 12 paths. foma[1]: words unbig unbiger unbigest unclear unclearer unclearest



#### Visualizing automata

- It can be useful to visualize FSAs as graphs
- Generic graph visualization software:
  - <u>http://www.graphviz.org/</u>
  - Defines a format for graphs
  - Graphs are rendered using layouting algorithms
  - Several options come with Graphviz, notably **dot**
- You also need a viewer:
  - A simple cross-platform option: ZGR Viewer
  - <u>http://zvtm.sourceforge.net/zgrviewer.html</u>



#### Visualizing automata

• Example: this little NLP logo:

digraph finite\_state\_machine {
 rankdir=LR;
 size="8,5"

node [shape = circle];

N -> L [ label = "t+1" ];

L -> P [ label = "t+2" ];

P-> P [ label = "self" ];

node [shape = doublecircle]; N P;



```
N t+1 t t+2 t
```

#### Our syllable example

- Foma [1]: print dot > some\_file\_name.dot
- Open with a GraphViz viewer (e.g. ZGRViewer)
  - Note the states and transitions
  - Start and end symbols (0 and double circle)



# Script files

- More often we'll define an automaton in a separate file
- Let's look at a script describing English numerals:
  - Download this script by Lauri Karttunen:

http://corpling.uis.georgetown.edu/amir/public/numeral.script



- Suppose we want to model the grammar of numbers from 1 – 99
- First we need one nine:
  - # Excerpt from Karttunen (2004)
  - # List the numbers from one to nine
  - **define** *OneToNine* [{one} | {two} | {three} | {four} | {five} | {six} | {seven} | {eight} | {nine}];



- We need to deal with teens + multiples of ten
- What kinds of morphemes are there?
- Schematically:
  - thir + [{teen}|{ty}]
  - four + [{teen}|{ty}]
  - fif + [{teen}|{ty}]
  - ...

Ten, eleven, twelve and twenty are separate



#### # Rules for teens:

**define** *TeenTen* [{thir} | {fif} | {six} | {seven} | {eigh} | {nine}];

**define** *Teens* [{ten} | {eleven} | {twelve} | [*TeenTen* | {four}] {teen}];

# Now things that can be followed by -ty: define TenStem [TeenTen | {twen} | {for}];



# Finally, allow twenty one, twenty two...: define Tens [TenStem {ty} ({-} OneToNine)];

# Now all possible numbers are: define OneToNinetyNine [ OneToNine | Teens | Tens ];

# Push our automaton to the stack for use
regex OneToNinetyNine;
print random-words
exit



## Result

foma[1]: **random-words** [1] fifty-five

- [1] seven
- [2] fifty
- [1] forty
- [1] fifty-nine
- [1] eleven
- [1] twenty
- [2] thirty
- [1] ten
- [1] nineteen
- [1] two
- [1] nine
- [1] ninety-seven

## Running the script

- To run a script file from the terminal:
  - foma -l numeral.script



#### What next

- This has been a very shallow introduction: 'regex with morphemes'
- FSMs can go much further
- Usually:
  - Two way translation between forms and analyses
  - Formally, Finite-State Transducers (FSTs)



#### What next

 More realistic examples: 398 bytes. 7 states, 7 arcs, 2 paths. foma[1]: up apply up> cats cat+N+Pl foma[1]: down apply down> cat+N+Pl cats

 The basis for generating inflected forms in NLG, analysis in morphologically rich NLU



#### Exercise for home – Japanese verbs

- You can practice writing finite-state morphologies on a language you are less familiar with
- Try modeling four verbs from the two major conjugation classes in Japanese:
  - -eru/-iru verbs:

• -u verbs:

*taberu* 'eat'*, nobiru* 'stretch' *yomu* 'read'*, kaku* 'write'



#### Exercise for home – Japanese verbs

- We can model the causative and passive inflections:
  - -iru/-eru verbs:
    - Drop 'ru'
    - Add saseru (causative) or rareru (passive)
    - or both: saserareru (be made to do something)
    - tabesaseru: make someone eat; nobirareru: be stretcheu
  - -u verbs:
    - Drop 'u'
    - Add aseru (causative) or areru (passive)
    - or both: aserareru
    - yomaserareru: be made to read





#### Exercise for home – Japanese verbs

- Produce a script that:
  - Defines the verb stems in each class
  - Defines the necessary suffixes
  - Combines the suffixes correctly with each class
- Using the words command, you should get all 3 possible inflected forms for all 4 verbs (12 forms):
  - tabesaseru, taberareu, tabesaserareru (be made to eat)

 If you want to learn more, send me an e-mail and stop by LING-362 in two weeks!

