



Towards a Data-centric View of Cloud Security

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Introduction

■ **Success of cloud**

- Economics of outsourcing data and computation
- Continued migration of applications to the cloud
- Amazon EC2, salesforce, Microsoft Office...

■ **Security: one barrier that prevents further success**

- Enforce the privacy and integrity of user data
- Current solutions mostly focus on OS and virtual machines

■ **Cloud applications are increasingly **interdependent****

Motivating Examples

- **Interconnecting enables more adaptable systems**
- **Online market-places**
 - Retail portals such as Yahoo!, Amazon serve as storefronts
 - Collect product and inventory information from sellers
 - Should prevent from ...
 - merchants querying each others' inventories and prices
 - communicating payment info with unauthorized parties
- **Social network services, outsourced data storage...**

Overview

- **A comprehensive solution should ...**

- go beyond OS and VM-centric security solutions
- securely share, verify, and trace data between applications

- **DS2 (Declarative Secure Distributed Systems)**

<http://netdb.cis.upenn.edu/ds2/>

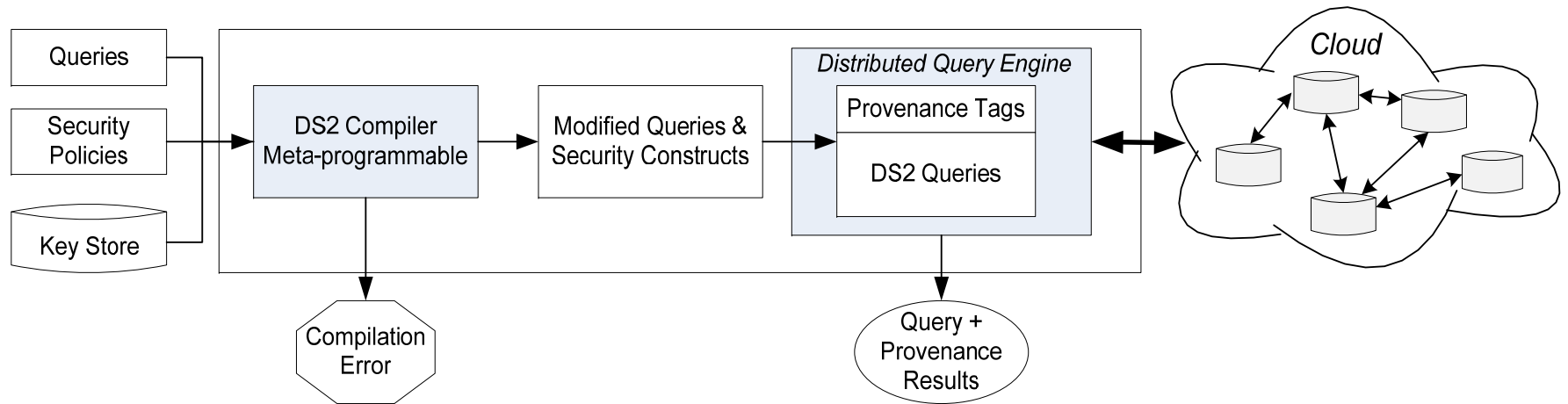
- Secure querying processing
- Declarative access control policies for data sharing
- System analysis and forensics using distributed provenance
- End-to-end verification of data partitioned across users



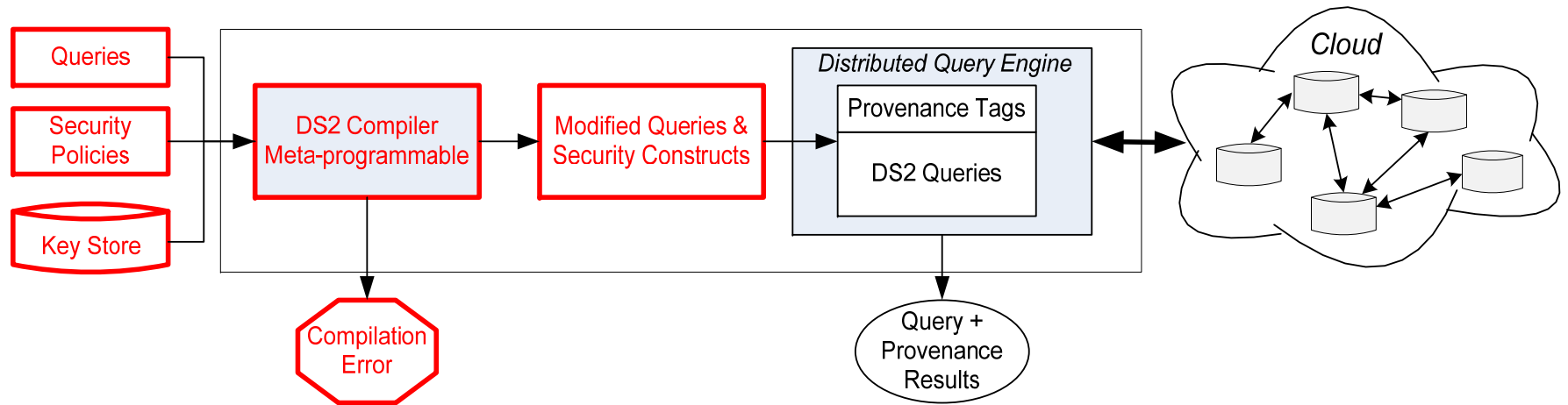
Outline

- Introduction
- Motivation
- DS2 Platform
 - Secure Data Processing
 - Declarative Access Control
 - Distributed Provenance
 - End-to-end query verification
- Conclusion

DS2 Platform

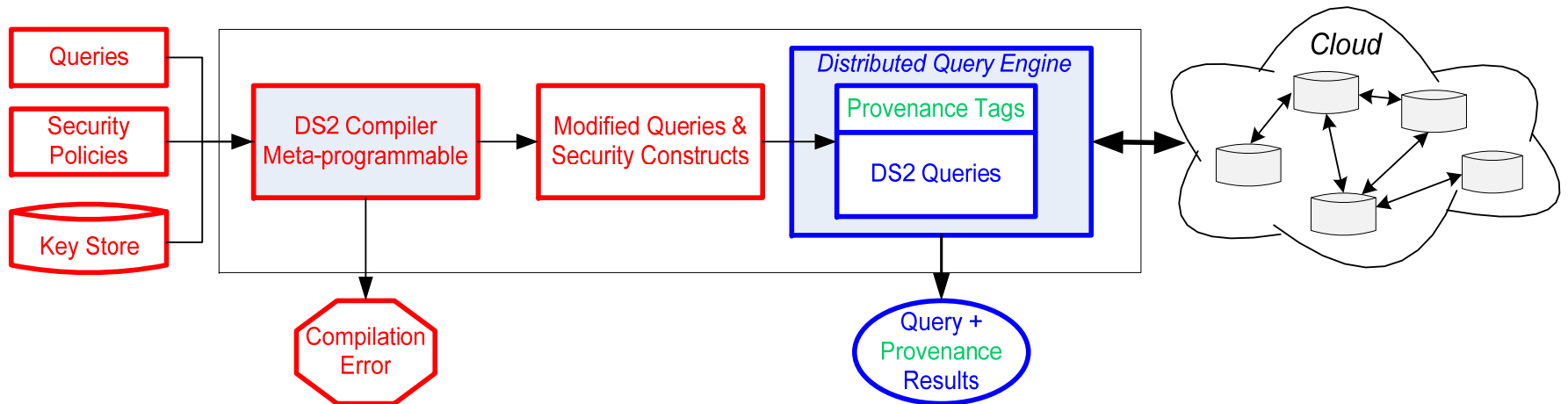


DS2 Platform



- Integration of access control policies
 - Meta-programmability

DS2 Platform



- Integration of access control policies
 - Meta-programmability
- Provenance-aware secure query processing
- End-to-end verification



Secure Query Processing

Zhou et al. **Unified Declarative Platform for Secure Networked Information Systems**, ICDE09

- **Compact specification of network protocols**
- ***Secure Network Datalog (SeNDlog)***
 - A distributed variant of *Datalog*
 - Continuous recursive queries over network state
 - Security Primitives
 - Rules within a *context*
 - Authenticated communication
- **A variety of secure distributed systems**
 - Secure network routing (S-BGP), DHTs, p2p query processing

Example: Authenticated Map-Reduce

At MW:

```
m1 map(ID,Content) :- file(MW,ID,Content).
```

```
m2 MW says emits(MW,Word,Num,Offset)@RW :-  
    word(Word,Num,Offset),  
    reduceWorker(RID,RW), RID=f_SHA1(Word).
```

■ In the context of Map Worker

- m1: Perform map operation on each file
 - m2: For each word in the document, pass it to the reducer according to the mapper-reducer mapping.
- ## ■ Authenticate outgoing tuples by tagging signatures

Example: Authenticated Map-Reduce

At RW:

r1 reduceTuple(Word,a_LIST<Num>) :-

MW says emits(MW,Word,Num,Offset).

r2 reduce(Word,List) :- reduceTuple(Word,List),

Master says rBegin(RW).

■ In the context of Reduce Worker

- r1: Group the received words, and maintain a list for each word
- r2: Perform reduce operation once received signal from Master

■ Verify the signatures of the incoming tuples

Example: Authenticated Map-Reduce

At RW:

r1 reduceTuple(Word,a_LIST<Num>) :-

MW says emits(MW,Word,Num,Offset).

Unified platform: protocol specs & security enforcement

Building blocks for more complex security policies

- r1: Group the received words, and maintain a list for which word
- r2: Perform reduce operation once received signal from Master
- **Verify the signatures of the incoming tuples**

Access Control

Marczak et al. SecureBlox: Customizable Secure Distributed Data Processing, SIGMOD10

■ View-based Access Control

- Horizontal and vertical partition of relational table
- Authorization + authentication
- Access ONLY to the secure views
- **How can we enforce this?**

At alice:

```
sv1 sview(Name,Dept) :- employee(Name,Dept,Salary), Salary < 5K.  
sv2 predsecview("employee","sview",U) :- authority says good(U).  
sv3 ret(Name,Dept)@U :- U says query("sview"), sview(Name, Dept).
```



Access Control

Marczak et al. SecureBlox: Customizable Secure Distributed Data Processing, SIGMOD10

■ Enforcement: meta-constraints

- Meta-model – rules as data
- Check the query format against schema constraints
- `says(U,R), body(R,A), functor(A,P) -> predsecview(_,P,U)`

■ Code Generation

- Automatic rewrite of queries to refer to security views
- Updates in the meta-model
- Customizable security constructs according to policy changes



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Distributed Provenance

Zhou et al. Efficient Querying and Maintenance of Network Provenance at Internet-Scale, SIGMOD10

■ Distributed provenance (or lineage)

- Explains the existence and derivation of any network state
- Maps naturally into various applications

■ Applications in cloud

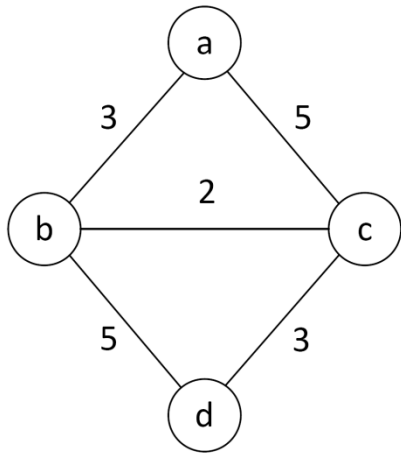
- Error detection, diagnosis, and forensics
- Mitigation: propagating corrections only to affected applications
- History-based trust management

Distributed Provenance

Zhou et al. Efficient Querying and Maintenance of Network Provenance at Internet-Scale, SIGMOD10

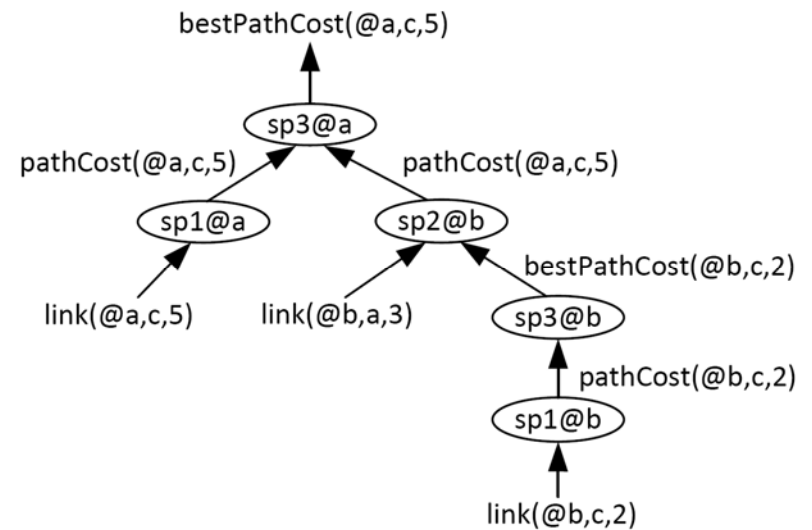
■ Data model – a directed graph

- Tuple and rule execution vertices
- Edges represent dataflows



link (@src, dest, cost)

src	dest	cost
a	b	3
a	c	5
b	a	3
b	c	2
b	d	5
c	a	5
c	b	2
c	d	3
d	b	5
d	c	3





Distributed Provenance

Zhou et al. Efficient Querying and Maintenance of Network Provenance at Internet-Scale, SIGMOD10

- **Data model – a directed graph**
 - Tuple and rule execution vertices
 - Edges represent dataflows
- **Maintenance and querying**
 - Maintained as distributed relational tables
 - Views of base and derived tuples
 - Querying performed as graph traversal
- **Reasonable overhead for distributed provenance**



End-to-end Query Verification

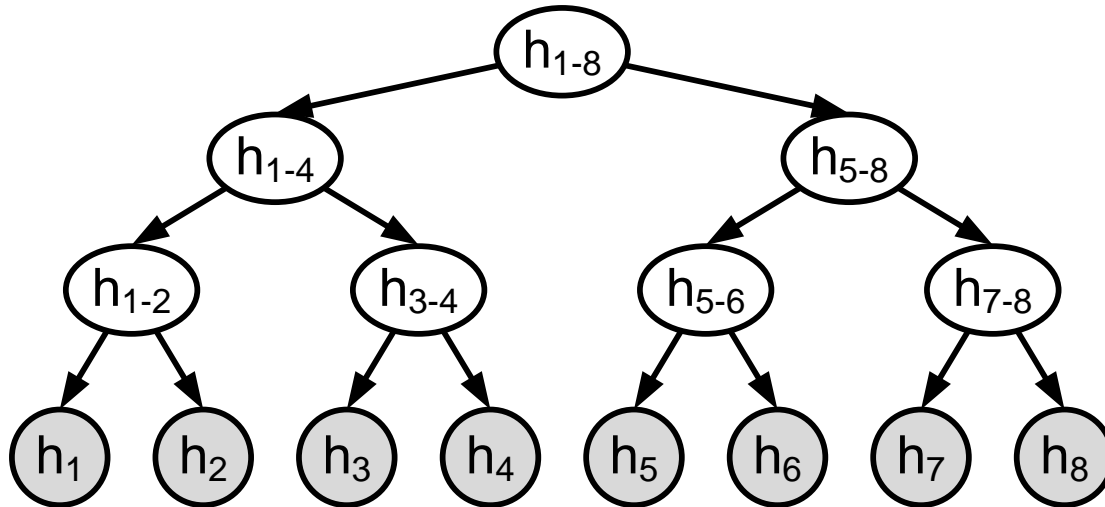
■ Threat Model

- The owner of the data is trustworthy
- Some fraction of the cloud that host the data could be malicious

■ Verification with MHT (Merkle Hash Tree)

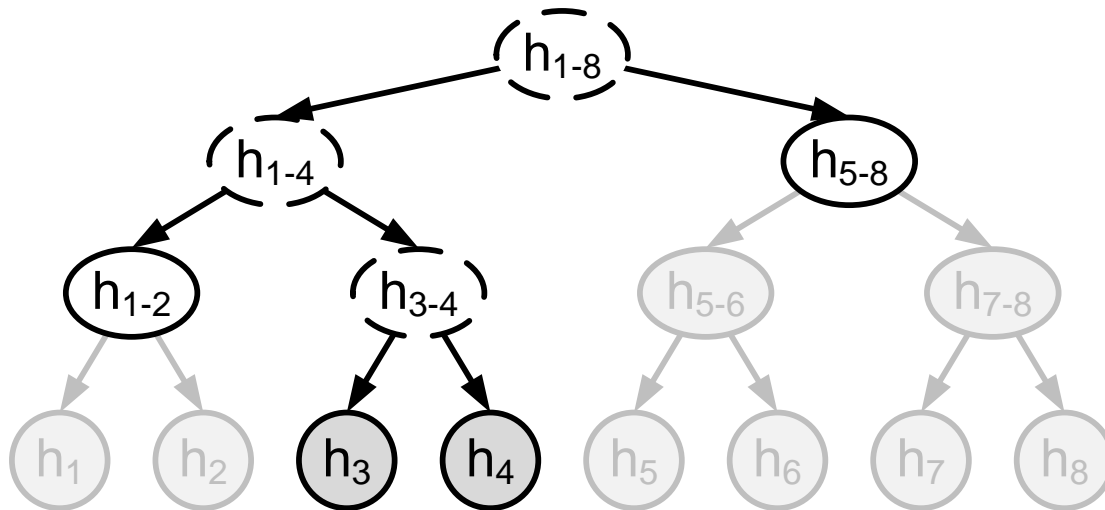
- Previously used to check correctness of outsourced databases
- Maintain hash hierarchy (MHT) on pre-sorted data
- VOs (verification objects) attached to query results
 - Signature over the root of MHT
 - Hash values required for re-computing the root of MHT

MHT Example



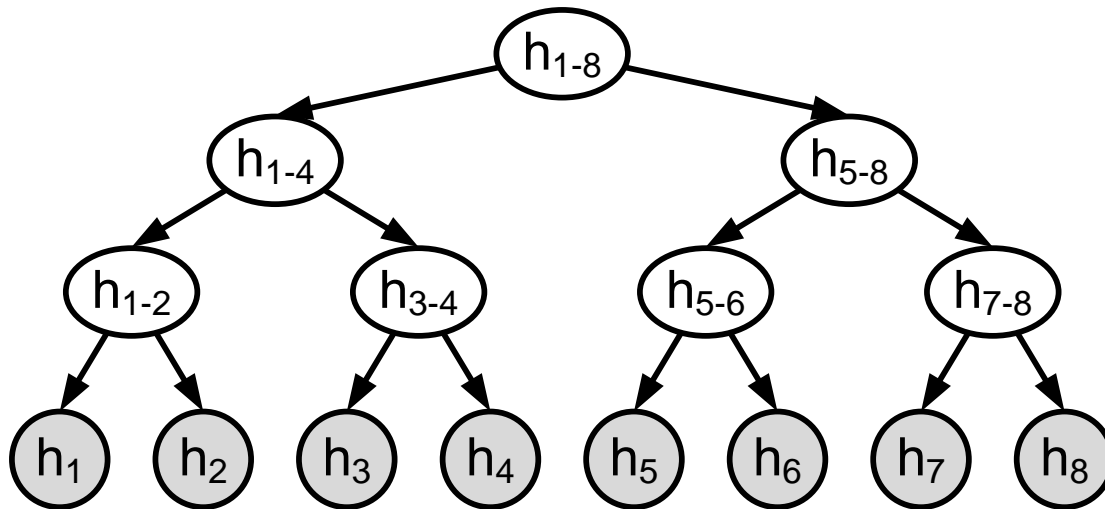
- Query Result = $\{x3\}$

MHT Example



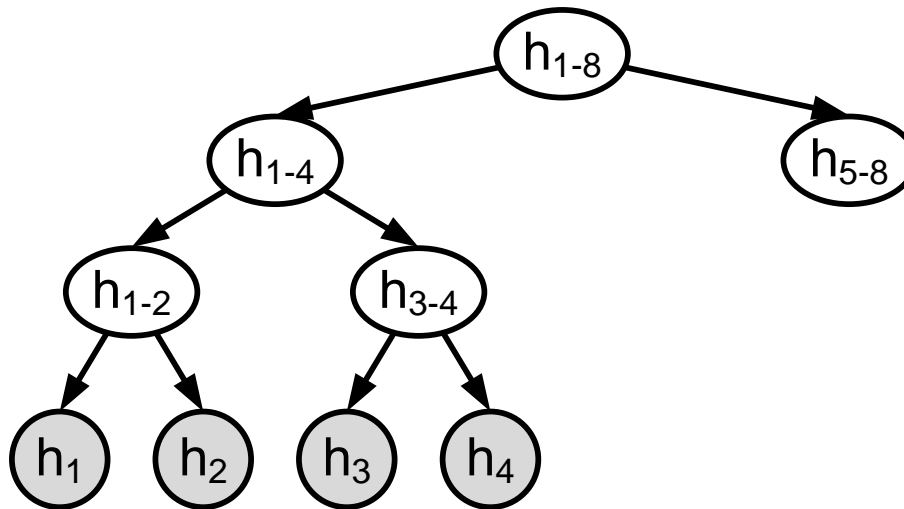
- Query Result = {x3}
- VO = {SIG(h₁₋₈), h₄, h₁₋₂, h₅₋₈}
- $\text{hash}(x_3) \mid h_4 \mid h_{1-2} \mid h_{5-8} == h_{1-8}?$

P-MHT Example



- Table is partitioned across three nodes
 - $X1 = \{x1, x2, x3\}$, $X2 = \{x4, x5, x6\}$, $X3 = \{x7, x8\}$
- Each node maintain a portion of MHT
 - Sufficient to generate the VOs for the tuples located on the node.

P-MHT Example



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Conclusion and Future Work

- Data-centric: go beyond OS/VM-centric solutions
- Security challenges faced by data-centric cloud security
 - Secure query processing and data sharing
 - Analysis and tracing of data flowing across applications
 - End-to-end verification
- Preliminary design of the DS2 Platform
- Future work
 - Close integration with cloud applications
 - Security guarantees for distributed provenance



Thank You...

Towards a Data-centric View of Cloud Security

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